The number of numerical semigroups of a given genus

Sergi Elizalde

Dartmouth College

SIAM Conference on Discrete Mathematics Minisymposium on Enumerative Combinatorics Austin, June 2010



The coin problem

Given coins of denominations c_1, c_2, \ldots, c_m ,

- what is the largest amount that cannot be obtained? (Frobenius problem)
- how many positive amounts cannot be obtained?

The coin problem

Given coins of denominations c_1, c_2, \ldots, c_m ,

- what is the largest amount that cannot be obtained? (Frobenius problem)
- how many positive amounts cannot be obtained?

Example:

With two coins of denominations 3 and 5, one can obtain

$$0, 3, 5, 6, 8, 9, 10, \dots$$

The coin problem

Given coins of denominations c_1, c_2, \ldots, c_m ,

- what is the largest amount that cannot be obtained? (Frobenius problem)
- how many positive amounts cannot be obtained?

Example:

With two coins of denominations 3 and 5, one can obtain

$$0, 3, 5, 6, 8, 9, 10, \dots$$

Such a set is called a *numerical semigroup*.



Definitions

A numerical semigroup is a set $\Lambda \subseteq \mathbb{N}_0 = \{0, 1, 2, \dots\}$ satisfying:

- ightharpoonup $0 \in \Lambda$,
- Λ is closed under addition,
- ▶ $\mathbb{N}_0 \setminus \Lambda$ is finite.

Definitions

A numerical semigroup is a set $\Lambda \subseteq \mathbb{N}_0 = \{0, 1, 2, \dots\}$ satisfying:

- 0 ∈ Λ,
- Λ is closed under addition,
- ▶ $\mathbb{N}_0 \setminus \Lambda$ is finite.

The elements in $\mathbb{N}_0 \setminus \Lambda$ are called *gaps*.

The *genus* of Λ is the number of gaps, denoted g.

The *Frobenius number* of Λ is the largest gap, denoted f.

Definitions

A numerical semigroup is a set $\Lambda \subseteq \mathbb{N}_0 = \{0, 1, 2, \dots\}$ satisfying:

- 0 ∈ Λ,
- Λ is closed under addition,
- ▶ $\mathbb{N}_0 \setminus \Lambda$ is finite.

The elements in $\mathbb{N}_0 \setminus \Lambda$ are called *gaps*.

The *genus* of Λ is the number of gaps, denoted g.

The *Frobenius number* of Λ is the largest gap, denoted f.

Example:

$$\Lambda = \{0, 4, 6, 8, 9, 10, 11, \dots\}$$
 $f = 7, g = 5$



Let n_g be the number of numerical semigroups of genus g.

$$g = 1: \quad \{0, 2, 3, 4, \dots\}$$

$$g = 2: \quad \{0, 2, 4, 5, 6, \dots\} \quad \{0, 3, 4, 5, 6, \dots\}$$

$$g = 3: \quad \{0, 2, 4, 6, 7, 8, \dots\} \quad \{0, 3, 4, 6, 7, 8, \dots\}$$

$$\{0, 3, 5, 6, 7, 8, \dots\} \quad \{0, 4, 5, 6, 7, 8, \dots\}$$

Every numerical semigroup Λ with $g \geq 1$ has a unique minimal set of generators $\mu_1, \mu_2, \dots, \mu_m$.

Every numerical semigroup Λ with $g \geq 1$ has a unique minimal set of generators $\mu_1, \mu_2, \dots, \mu_m$.

If
$$\mu_1 < \dots < \mu_r < f < \underbrace{\mu_{r+1} < \dots < \mu_m}_{\textit{effective generators}}$$
 , we write

$$\Lambda = \langle \mu_1, \dots, \mu_r | \mu_{r+1}, \dots, \mu_m \rangle.$$

Every numerical semigroup Λ with $g \geq 1$ has a unique minimal set of generators $\mu_1, \mu_2, \dots, \mu_m$.

If
$$\mu_1 < \dots < \mu_r < f < \underbrace{\mu_{r+1} < \dots < \mu_m}_{effective\ generators}$$
 , we write

$$\Lambda = \langle \mu_1, \dots, \mu_r | \mu_{r+1}, \dots, \mu_m \rangle.$$

 $\Lambda \cup \{f\}$ is a numerical semigroup of genus g-1.

Every numerical semigroup Λ with $g \geq 1$ has a unique minimal set of generators $\mu_1, \mu_2, \dots, \mu_m$.

If
$$\mu_1 < \dots < \mu_r < f < \underbrace{\mu_{r+1} < \dots < \mu_m}_{effective\ generators}$$
 , we write

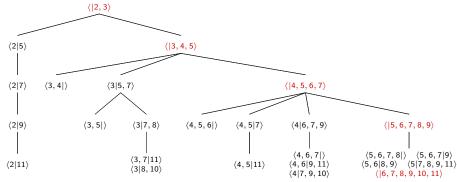
$$\Lambda = \langle \mu_1, \dots, \mu_r | \mu_{r+1}, \dots, \mu_m \rangle.$$

 $\Lambda \cup \{f\}$ is a numerical semigroup of genus g-1.

Example:
$$\{0, 4, 6, 8, 9, 10, 11, \dots\} = \langle 4, 6 | 9, 11 \rangle$$
, $\{0, 4, 6, 7, 8, 9, 10, 11, \dots\} = \langle 4 | 6, 7, 9 \rangle$



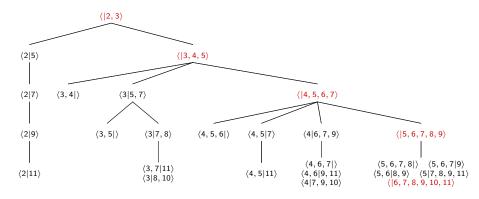
The tree \mathcal{T} of numerical semigroups



Consider the tree \mathcal{T} with root $\{0, 2, 3, 4, \dots\} = \langle |2, 3\rangle$ where

- ▶ the parent of each Λ is $\Lambda \cup \{f\}$,
- ▶ the children of each $\Lambda = \langle \mu_1, \dots, \mu_r | \mu_{r+1}, \dots, \mu_{r+e} \rangle$ are $\Lambda \setminus \{\mu_{r+i}\}$, with $1 \le i \le e$.

The tree \mathcal{T} of numerical semigroups



The number of nodes at level g is n_g . We will bound n_g by approximating this tree with simpler trees, keeping track of the number of effective generators of each node.



 $O_g = \{0, g+1, g+2, g+3, \dots\} = \langle |g+1, g+2, \dots, 2g+1 \rangle$ is the *ordinary semigroup* of genus g. It has g+1 effective generators.

$$O_g = \{0, g+1, g+2, g+3, \dots\} = \langle |g+1, g+2, \dots, 2g+1 \rangle$$
 is the *ordinary semigroup* of genus g . It has $g+1$ effective generators.

The g+1 children of O_g have $0, 1, 2, \dots, g-2, g, g+2$ effective generators respectively (the last child being O_{g+1}).

$$O_g = \{0, g+1, g+2, g+3, \dots\} = \langle |g+1, g+2, \dots, 2g+1 \rangle$$
 is the *ordinary semigroup* of genus g . It has $g+1$ effective generators.

The g+1 children of O_g have $0,1,2,\ldots,g-2,g,g+2$ effective generators respectively (the last child being O_{g+1}).

We write this as

$$\overline{(g+1)} \longrightarrow (0)(1)\dots(g-2)(g)\overline{(g+2)},$$

$$O_g = \{0, g+1, g+2, g+3, \dots\} = \langle |g+1, g+2, \dots, 2g+1 \rangle$$
 is the *ordinary semigroup* of genus g . It has $g+1$ effective generators.

The g+1 children of O_g have $0, 1, 2, \ldots, g-2, g, g+2$ effective generators respectively (the last child being O_{g+1}).

We write this as

$$\overline{(g+1)} \longrightarrow (0)(1)\dots(g-2)(g)\overline{(g+2)},$$

or equivalently as

$$\overline{(e)} \longrightarrow (0)(1)\dots(e-3)(e-1)\overline{(e+1)}.$$



Let $\Lambda = \langle \mu_1, \dots, \mu_r | \mu_{r+1}, \dots, \mu_{r+e} \rangle$ be a non-ordinary semigroup. Then, for $1 \le i \le e$,

$$\Lambda \setminus \{\mu_{r+i}\} = \begin{cases} \langle \mu_1, \dots, \mu_{r+i-1} | \underbrace{\mu_{r+i+1}, \dots, \mu_{r+e}} \rangle & \text{or} \\ e-i & \text{effective gen.} \\ \langle \mu_1, \dots, \mu_{r+i-1} | \underbrace{\mu_{r+i+1}, \dots, \mu_{r+e}, \mu_1 + \mu_{r+i}}_{e-i+1} \rangle. \end{cases}$$

Let $\Lambda = \langle \mu_1, \dots, \mu_r | \mu_{r+1}, \dots, \mu_{r+e} \rangle$ be a non-ordinary semigroup. Then, for $1 \le i \le e$,

$$\Lambda \setminus \{\mu_{r+i}\} = \begin{cases} \langle \mu_1, \dots, \mu_{r+i-1} | \underbrace{\mu_{r+i+1}, \dots, \mu_{r+e}} \rangle & \text{or} \\ e-i & \text{effective gen.} \\ \langle \mu_1, \dots, \mu_{r+i-1} | \underbrace{\mu_{r+i+1}, \dots, \mu_{r+e}, \mu_1 + \mu_{r+i}} \rangle. \\ e-i+1 & \text{effective gen.} \end{cases}$$

Ex: The children of $\langle 4|6,7,9 \rangle$ are $\langle 4,6,7| \rangle$, $\langle 4,6|9,11 \rangle$, $\langle 4|7,9,10 \rangle$.

Let $\Lambda = \langle \mu_1, \dots, \mu_r | \mu_{r+1}, \dots, \mu_{r+e} \rangle$ be a non-ordinary semigroup. Then, for $1 \le i \le e$,

$$\Lambda \setminus \{\mu_{r+i}\} = \begin{cases} \langle \mu_1, \dots, \mu_{r+i-1} | \underbrace{\mu_{r+i+1}, \dots, \mu_{r+e}} \rangle & \text{or} \\ e-i & \text{effective gen.} \\ \langle \mu_1, \dots, \mu_{r+i-1} | \underbrace{\mu_{r+i+1}, \dots, \mu_{r+e}, \mu_1 + \mu_{r+i}} \rangle. \\ e-i+1 & \text{effective gen.} \end{cases}$$

Ex: The children of $\langle 4|6,7,9\rangle$ are $\langle 4,6,7|\rangle$, $\langle 4,6|9,11\rangle$, $\langle 4|7,9,10\rangle$. (3) \longrightarrow (0)(2)(3)

Let $\Lambda = \langle \mu_1, \dots, \mu_r | \mu_{r+1}, \dots, \mu_{r+e} \rangle$ be a non-ordinary semigroup. Then, for $1 \le i \le e$,

$$\Lambda \setminus \{\mu_{r+i}\} = \begin{cases} \langle \mu_1, \dots, \mu_{r+i-1} | \underbrace{\mu_{r+i+1}, \dots, \mu_{r+e}} \rangle & \text{or} \\ e-i & \text{effective gen.} \\ \langle \mu_1, \dots, \mu_{r+i-1} | \underbrace{\mu_{r+i+1}, \dots, \mu_{r+e}, \mu_1 + \mu_{r+i}} \rangle. \\ e-i+1 & \text{effective gen.} \end{cases}$$

Ex: The children of $\langle 4|6,7,9 \rangle$ are $\langle 4,6,7| \rangle$, $\langle 4,6|9,11 \rangle$, $\langle 4|7,9,10 \rangle$.

$$(3) \longrightarrow (0)(2)(3)$$

In general, $(e) \longrightarrow (j_1)(j_2)\dots(j_e)$, where $j_i \in \{i-1,i\}$.

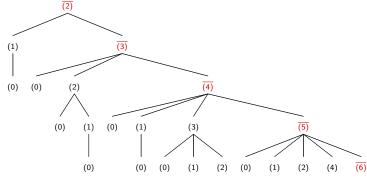


A lower bound

Consider the generating tree with root (2) and succession rules

$$(e) \longrightarrow (0)(1)\dots(e-3)(e-1)(e+1),$$

$$(e) \longrightarrow (0)(1)\dots(e-1).$$



A lower bound

Consider the generating tree with root (2) and succession rules

$$(e) \longrightarrow (0)(1)\dots(e-3)(e-1)(e+1),$$

$$(e) \longrightarrow (0)(1)\dots(e-1).$$

This tree can be embedded in \mathcal{T} , so its number of nodes at level g is a lower bound on n_g .

A lower bound

Consider the generating tree with root (2) and succession rules

$$(e) \longrightarrow (0)(1)\dots(e-3)(e-1)(e+1),$$

$$(e) \longrightarrow (0)(1)\dots(e-1).$$

This tree can be embedded in \mathcal{T} , so its number of nodes at level g is a lower bound on n_g .

From the succession rules, the generating function for the number of nodes at each level is

$$\frac{t(1+t+t^2)}{1-t-t^2}=t+2t^2+4t^3+6t^4+10t^5+\cdots=t+\sum_{g>2}2F_g\,t^g.$$

So, for $g \ge 2$,

$$n_g \geq 2F_g$$
.

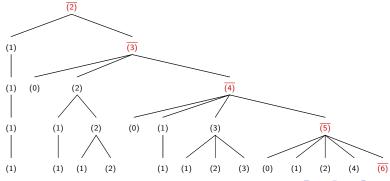


An upper bound

Consider the generating tree with root (2) and succession rules

$$(e) \longrightarrow (0)(1)\dots(e-3)(e-1)\overline{(e+1)},$$

$$(e) \longrightarrow (1)(2)\dots(e).$$



An upper bound

Consider the generating tree with root (2) and succession rules

$$(e) \longrightarrow (0)(1)\dots(e-3)(e-1)(e+1),$$

$$(e) \longrightarrow (1)(2)\dots(e).$$

 \mathcal{T} can be embedded in this tree, so its number of nodes at level g is an upper bound on n_g .

An upper bound

Consider the generating tree with root (2) and succession rules

$$\overline{(e)} \longrightarrow (0)(1)\dots(e-3)(e-1)\overline{(e+1)},$$

$$(e) \longrightarrow (1)(2)\dots(e).$$

 \mathcal{T} can be embedded in this tree, so its number of nodes at level g is an upper bound on n_g .

From the succession rules, the generating function for the number of nodes at each level is

$$\frac{t(1-t-t^3)}{(1-t)(1-2t)}=t+2t^2+4t^3+7t^4+13t^5+\ldots$$

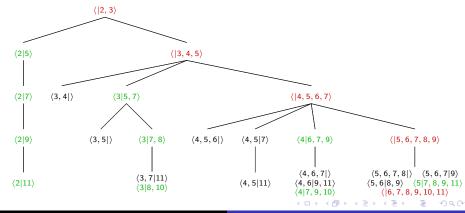
So, for $g \geq 3$,

$$n_g \leq 1 + 3 \cdot 2^{g-3}.$$

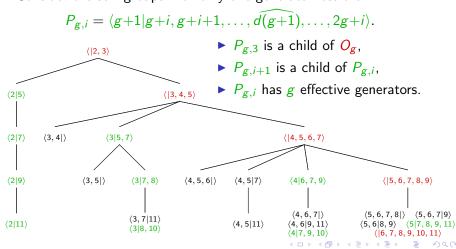


Consider the semigroups with only one generator less than f:

$$P_{g,i} = \langle g+1|g+i,g+i+1,\ldots,\widehat{d(g+1)},\ldots,2g+i\rangle.$$



Consider the semigroups with only one generator less than f:



Consider the semigroups with only one generator less than f:

$$P_{g,i} = \langle g+1|g+i, g+i+1, \dots, \widehat{d(g+1)}, \dots, 2g+i \rangle.$$

- $\triangleright P_{g,3}$ is a child of O_g ,
- $ightharpoonup P_{g,i+1}$ is a child of $P_{g,i}$,
- $ightharpoonup P_{g,i}$ has g effective generators.

We write
$$(\widetilde{g}) \longrightarrow (j_1)(j_2) \dots (j_{g-1})(\widetilde{g})$$
 where $j_i \in \{i-1,i\}$.

Consider the semigroups with only one generator less than f:

$$P_{g,i} = \langle g+1|g+i, g+i+1, \dots, \widehat{d(g+1)}, \dots, 2g+i \rangle.$$

- $ightharpoonup P_{g,3}$ is a child of O_g ,
- $ightharpoonup P_{g,i+1}$ is a child of $P_{g,i}$,
- $ightharpoonup P_{g,i}$ has g effective generators.

We write
$$(\widetilde{g}) \longrightarrow (j_1)(j_2) \dots (j_{g-1})(\widetilde{g})$$
 where $j_i \in \{i-1, i\}$.

The succession rules for the new tree are

$$(e) \longrightarrow (0)(1)\dots(e-3)(e-1)(e+1),$$

$$(e) \longrightarrow (0)(1)\dots(e-2)(e),$$

$$(e) \longrightarrow (0)(1)\dots(e-1).$$

Consider the semigroups with only one generator less than f:

$$P_{g,i} = \langle g+1|g+i, g+i+1, \dots, \widehat{d(g+1)}, \dots, 2g+i \rangle.$$

- $ightharpoonup P_{g,3}$ is a child of O_g ,
- $ightharpoonup P_{g,i+1}$ is a child of $P_{g,i}$,
- $ightharpoonup P_{g,i}$ has g effective generators.

We write
$$(\widetilde{g}) \longrightarrow (j_1)(j_2) \dots (j_{g-1})(\widetilde{g})$$
 where $j_i \in \{i-1, i\}$.

The succession rules for the new tree are

$$(e) \longrightarrow (0)(1)\dots(e-3)(e-1)(e+1),$$

$$(e) \longrightarrow (0)(1)\dots(e-2)(e),$$

$$(e) \longrightarrow (0)(1)\dots(e-1).$$

Counting the nodes gives an improved lower bound:

$$n_g \geq F_{g+2} - 1 \geq 2F_g.$$



Idea: Use a second label to keep track of the number of strong generators of each semigroup. An effective gen. $\mu \in \Lambda$ is called *strong* if $\mu + \mu_1$ is a generator of $\Lambda \setminus \{\mu\}$.

Idea: Use a second label to keep track of the number of strong generators of each semigroup. An effective gen. $\mu \in \Lambda$ is called *strong* if $\mu + \mu_1$ is a generator of $\Lambda \setminus \{\mu\}$.

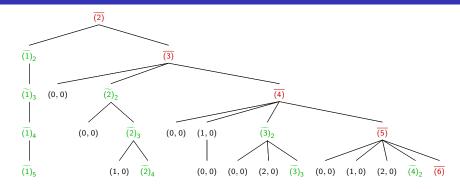
We bound the number of strong gen. in terms on the number of strong gen. of the parent. The succession rules become

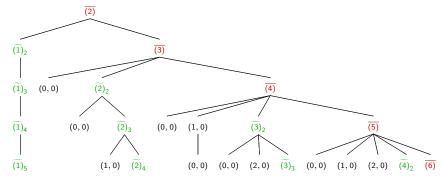
$$\begin{array}{ccc} (e) & \longrightarrow & (0,0)(1,0)\dots(e-3,0)(e-1)_2(e+1), \\ \widetilde{(e)}_k & \longrightarrow & (0,0)(1,0)\dots(e-\sigma-1,0)(e-\sigma+1,0)(e-\sigma+2,1)\dots(e-1,\sigma-2)(\widetilde{e)}_{k+1}, \end{array}$$

$$(e,s) \longrightarrow (0,0)(1,0)\dots(e-s-1,0)(e-s+1,0)(e-s+2,1)\dots(e,s-1).$$

where

$$\sigma = \sigma(e,k) := \begin{cases} k & \text{if } 2 \leq k \leq \lceil e/2 \rceil, \\ k-1 & \text{if } \lceil e/2 \rceil < k \leq e, \quad (\# \text{ of strong gen. of } P_{e,k+1}) \\ e & \text{if } k > e. \end{cases}$$





The coefficients of its corresponding generating function

$$\frac{t\left(1-t^2-2t^3-3t^4+t^5+2t^6+3t^7+3t^8+t^9\right)}{(1+t)(1-t)(1-t-t^2)(1-t-t^3)(1-t^3-2t^4-2t^5-t^6)}$$

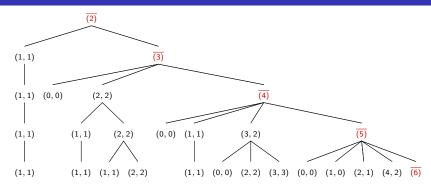
give a better lower bound on n_g .

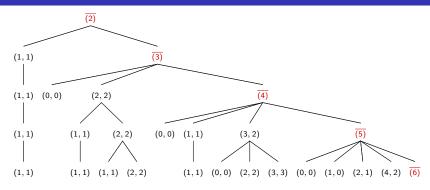


Idea: use a second label to keep track of the number of healthy generators of each semigroup. An effective gen. $\mu \in \Lambda$ is called healthy if $\mu + \mu_1 \leq 2g + 3$. Strong generators are always healthy.

Idea: use a second label to keep track of the number of healthy generators of each semigroup. An effective gen. $\mu \in \Lambda$ is called healthy if $\mu + \mu_1 \leq 2g + 3$. Strong generators are always healthy.

We bound the number of healthy gen. in terms on the number of effective and healthy gen. of the parent. The succession rules become





The coefficients of its corresponding generating function

$$t\,\frac{2-3t+t^2-4t^3+3t^4-2t^5+t(1-t-t^3)\sqrt{(1+2t)/(1-2t)}}{2(1-3t+3t^2-3t^3+4t^4-3t^5+2t^6)}$$

give the best known upper bound on n_g .



Numerical semigroups Easy bounds on n_g Improved bounds on n_g Better lower bounds A better upper bound Table of bounds

Г	g	2F _g	$F_{g+2} - 1$	lower bound	ng	upper bound	$1 + 3 \cdot 2^{g-3}$
	1		1	1	1	1	
	2	2	2	2	2	2	
	3	4	4	4	4	4	4
	4	6	7	7	7	7	7
	5	10	12	12	12	13	13
	6	16	20	22	23	24	25
	7	26	33	37	39	44	49
	8	42	54	62	67	81	97
	9	68	88	104	118	151	193
	10	110	143	175	204	280	385
	11	178	232	291	343	525	769
	12	288	376	482	592	984	1537
	13	466	609	796	1001	1859	3073
	14	754	986	1315	1693	3511	6145
	15	1220	1596	2166	2857	6682	12289
	16	1974	2583	3559	4806	12709	24577
	17	3194	4180	5838	8045	24334	49153
	18	5168	6764	9569	13467	46565	98305
	19	8362	10945	15665	22464	89626	196609
	20	13530	17710	25612	37396	172381	393217
	21	21892	28656	41831	62194	333262	786433
	22	35422	46367	68270	103246	643733	1572865
	23	57314	75024	111337	170963	1249147	3145729
	24	92736	121392	181438	282828	2421592	6291457
	25	150050	196417	295480	467224	4713715	12582913
	26	242786	317810	480938	770832	9165792	25165825
	27	392836	514228	782408	1270267	17888456	50331649
	28	635622	832039	1272250	2091030	34873456	100663297
	29	1028458	1346268	2067870	3437839	68212220	201326593
	30	1664080	2178308	3359757	5646773	133269997	402653185
	31	2692538	3524577	5456862	9266788	261167821	805306369
	32	4356618	5702886	8860132	15195070	511211652	1610612737



Open problems

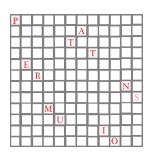
lacksquare $\lim_{g o\infty}rac{n_{g+1}}{n_g}=rac{1+\sqrt{5}}{2}$ (conjectured by Maria Bras-Amorós),

Open problems

- $ightharpoonup \ \lim_{g o\infty}rac{n_{g+1}}{n_g}=rac{1+\sqrt{5}}{2}$ (conjectured by Maria Bras-Amorós),
- ▶ $n_{g+1} \ge n_g$ for all g,

Open problems

- $ightharpoonup \ \lim_{g o\infty}rac{n_{g+1}}{n_{\sigma}}=rac{1+\sqrt{5}}{2}$ (conjectured by Maria Bras-Amorós),
- ▶ $n_{g+1} \ge n_g$ for all g,



Eighth International Conference on **Permutation Patterns**, *PP 2010*

August 9-13, Dartmouth College, Hanover, NH

Invited speakers:

- ▶ Nik Ruškuc, University of St Andrews
- Richard Stanley, MIT

 $\verb|http://math.dartmouth.edu/\sim|pp2010|$